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(* apl-2013-06-24c.v
 * Programmer: Mayer Goldberg, 2013
 *)
Require Import Arith.
Theorem our_plus_0_l: \forall a, 0 + a = a.
Proof.
  intro a.
  unfold plus; fold plus; reflexivity.
Qed.
Theorem our_plus_0r: \forall a, a + 0 = a.
Proof.
  induction a.
  reflexivity.
  unfold plus; fold plus.
  rewrite IHa; reflexivity.
Qed.
Lemma L1: \forall a b, S a + b = S (a + b).
Proof.
  fold plus; unfold plus; reflexivity.
Qed.
Lemma L2: \forall a b, a + S b = S (a + b).
Proof.
  induction a.
  induction b.
  reflexivity.
  repeat rewrite our_plus_0_l; reflexivity.
  induction b.
  rewrite our_plus_0_r; rewrite L1.
  rewrite IHa; rewrite our_plus_0_r; reflexivity.
  rewrite L1; rewrite IHa; rewrite L1; reflexivity.
Qed.
Theorem our_plus_comm: \forall a \ b, \ a + b = b + a.
Proof.
  induction a.
  intro b.
  rewrite our_plus_0_l, our_plus_0_r; reflexivity.
  induction b.
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rewrite our_plus_0_l; rewrite our_plus_0_r; reflexivity.
  repeat rewrite L1.
  repeat rewrite L2.
  rewrite IHa.
  reflexivity.
Qed.
Theorem our_mult_0l: \forall a, 0 \times a = 0.
Proof.
  intro a.
  unfold mult; reflexivity.
Qed.
Theorem our_mult_0r: \forall a, a \times 0 = 0.
Proof.
  induction a.
  unfold mult; reflexivity.
  unfold mult; fold mult.
  rewrite our_plus_0_l.
  exact IHa.
Qed.
Theorem our_mult_1_l: \forall a, 1 \times a = a.
Proof.
  intro a.
  unfold mult.
  rewrite our_plus_0_r; reflexivity.
Qed.
Theorem our\_mult\_1\_r: \forall a, a \times 1 = a.
Proof.
  induction a.
  rewrite our_mult_0_l; reflexivity.
  unfold mult; fold mult.
  unfold plus.
  rewrite IHa.
  reflexivity.
Qed.
Theorem our_plus_assoc: \forall a \ b \ c, \ a + (b + c) = a + b + c.
Proof.
  induction a, b, c.
  reflexivity.
  repeat rewrite our_plus_0_l; reflexivity.
  repeat rewrite our_plus_0_l; repeat rewrite our_plus_0_r; reflexivity.
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repeat rewrite our_plus_0_l; reflexivity.
  repeat rewrite our_plus_0r_; reflexivity.
  rewrite our_plus_0_l, our_plus_0_r; reflexivity.
  repeat rewrite our_plus_0 r; reflexivity.
  repeat rewrite L1; repeat rewrite L2; rewrite L1; rewrite IHa; reflexivity.
Qed.
Lemma L3: \forall a b, S a \times b = b + a \times b.
Proof.
  intros a b.
  unfold mult; fold mult; reflexivity.
Qed.
Lemma L4: \forall a b, a \times S b = a + a \times b.
Proof.
  induction a.
  intro b.
  repeat rewrite our_mult_0_l.
  rewrite our_plus_0_l; reflexivity.
  induction b.
  rewrite our_mult_1_r, our_mult_0_r, our_plus_0_r; reflexivity.
  repeat rewrite L3.
  repeat rewrite IHa.
  repeat rewrite our_plus_assoc.
  replace (S (S b) + a) with (S a + S b).
  reflexivity.
  repeat rewrite L1.
  rewrite L2, our_plus_comm; reflexivity.
Qed.
Theorem our_mult_comm: \forall a b, a \times b = b \times a.
Proof.
  induction a.
  intro b.
  rewrite our_mult_0_l, our_mult_0_r; reflexivity.
  induction b.
  rewrite our_mult_0_l, our_mult_0_r; reflexivity.
  repeat rewrite L3.
  repeat rewrite L4.
  repeat rewrite our_plus_assoc.
  \texttt{rewrite} \leftarrow IHa.
  repeat rewrite L1.
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replace (b + a) with (a + b).
  reflexivity.
  apply our_plus_comm.
Qed.
Theorem our\_mult\_plus\_distr\_r: \forall a \ b \ c, \ (a + b) \times c = a \times c + b \times c.
Proof.
  induction a, b, c.
  repeat rewrite our_mult_0r; reflexivity.
  repeat rewrite our_mult_0_l; reflexivity.
  repeat rewrite our_mult_0_r; reflexivity.
  repeat rewrite our_mult_0_l; repeat rewrite our_plus_0_l; reflexivity.
  repeat rewrite our_mult_0_r; reflexivity.
  rewrite our_plus_0_r, our_mult_0_l, our_plus_0_r; reflexivity.
  repeat rewrite our_mult_0_r; reflexivity.
  rewrite L1.
  rewrite L2.
  repeat rewrite L3.
  repeat rewrite L_4.
  rewrite IHa.
  repeat rewrite our_plus_assoc.
  repeat rewrite L1.
  repeat rewrite L2.
  repeat rewrite L1.
  rewrite (our_plus_comm (c + a + a \times c) c).
  repeat rewrite our_plus_assoc.
  replace ((c + c + a) + b + (a \times c)) with ((c + c + a) + (a \times c) + b).
  reflexivity.
  repeat rewrite \leftarrow our_plus_assoc.
  rewrite (our_plus_comm (a \times c) b).
  reflexivity.
Qed.
Theorem our\_mult\_plus\_distr\_l: \forall a \ b \ c, \ a \times (b + c) = a \times b + a \times c.
Proof.
  intros a b c.
  rewrite our_mult_comm.
  rewrite our_mult_plus_distr_r.
  rewrite our_mult_comm at 1.
  rewrite (our_mult_comm c a).
  reflexivity.
Qed.
Lemma acPbc: \forall a \ b \ c, \ a = b \rightarrow a + c = b + c.
Proof.
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intros a b c H.
  rewrite H; reflexivity.
Qed.
Theorem our_mult_assoc: \forall a \ b \ c, \ a \times (b \times c) = a \times b \times c.
Proof.
  induction a, b, c.
  repeat rewrite our_mult_0r; reflexivity.
  repeat rewrite our_mult_0_l, our_mult_0_r; reflexivity.
  repeat rewrite our_mult_0_l, our_mult_0_r; reflexivity.
  repeat rewrite our_mult_0l; reflexivity.
  repeat rewrite our_mult_0r; reflexivity.
  rewrite our_mult_0_l, our_mult_0_r; reflexivity.
  repeat rewrite our_mult_0_r; reflexivity.
  repeat rewrite L3, L4.
  repeat rewrite our_plus_assoc.
  repeat rewrite our_mult_plus_distr_r.
  repeat rewrite our_mult_plus_distr_l.
  repeat rewrite our_plus_assoc.
  rewrite IHa.
  rewrite (acPbc (S \ b + c + b \times c + a \times S \ b + a \times c) (S \ b + a + a \times b + S \ b \times c + a)
a \times c) (a \times b \times c)).
  reflexivity.
  rewrite (acPbc (S \ b + c + b \times c + a \times S \ b) (S \ b + a + a \times b + S \ b \times c) (a \times c)).
  reflexivity.
  rewrite L3.
  rewrite L4.
  repeat rewrite \leftarrow our_plus_assoc.
  repeat rewrite (our_plus_comm (S b) _).
  \texttt{rewrite} \; (acPbc \; (c + (b \times c + (a + a \times b))) \; (a + (a \times b + (c + b \times c))) \; (S \; b)).
  reflexivity.
  rewrite (our_plus_comm \ (b \times c) \ (a + a \times b)).
  repeat rewrite our_plus_assoc.
  rewrite (acPbc (c + a + a \times b) (a + a \times b + c) (b \times c)).
  reflexivity.
  rewrite \leftarrow (our_plus_assoc \ a \ (a \times b) \ c).
  rewrite (our_plus_comm (a \times b) c).
  rewrite our_plus_assoc.
  rewrite (our_plus_comm a c).
  reflexivity.
Qed.
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